V. V. V'yugin

Institute for Information Transmission Problems Russian Academy of Sciences

SAGA 2009, Sapporo 26-28 October 2009



・ ロ ・ ・ 雪 ・ ・ 目 ・

The talk in outline

The talk in outline

- Arbitrage strategies based on difference between macro and micro volatilities of a time series
- Prediction with expert advice Hannan's algorithm in case of unbounded gains
- Merging the arbitrage strategies using Hannan's algorithm

A D > A P > A D > A D >

Results of experiments

Macro and micro volatilities of a time series

Consider a discrete time series of stock prices

$$S_0 = S(0), S_1 = S(T/(KM)), S_2 = S(2T/(KM))..., S_{KM} = S(T).$$

The macro volatility is represented by the sum

$$\sum_{i=0}^{K-1} (S_{(i+1)T} - S_{iT})^2,$$

and the micro volatility is represented by the sum

$$\sum_{i=0}^{kT-1} (\Delta S_i)^2,$$

where $\Delta S_i = S_{i+1} - S_i$, i = 1, ..., KTIn this paper for simplicity we consider the case K = 1.



Stock price



Evolution of the price of a stock

◆□ > ◆□ > ◆臣 > ◆臣 >

Arbitrage strategies

The difference between macro and micro volatility sums

$$(S_T - S_0)^2 - \sum_{t=0}^{T-1} (\Delta S_t)^2 = (\sum_{t=0}^{T-1} \Delta S_t)^2 - \sum_{t=0}^{T-1} (\Delta S_t)^2 = \sum_{t=0}^{T-1} 2(S_t - S_0) \Delta S_t$$

Vovk, V.: A game-theoretic explanation of the \sqrt{dt} effect, (2003), http://www.probabilityandfinance.com

< ロ > < 回 > < 回 > < 回 > < 回 >

Arbitrage strategies

We consider two strategies

$$s_t^1 = 2(S_t - S_0)\Delta S_t,$$

 $s_t^2 = -2(S_t - S_0)\Delta S_t.$

These strategies earn the incomes on steps t = 1, ..., T

$$s_{1:T}^1 = \sum_{t=0}^{T-1} s_t^1 = (S_T - S_0)^2 - \sum_{t=0}^{T-1} (\Delta S_t)^2,$$

 $s_{1:T}^2 = \sum_{t=0}^{T-1} s_t^2 = \sum_{t=0}^{T-1} (\Delta S_t)^2 - (S_T - S_0)^2.$



Zero-sum game



Gains and losses of two experts.

Green line - micro volatility expert; Blue line - macro volatility expert.

• • • • • • • • • • • •

Prediction with expert advice

We use methods of prediction with expert advice to merge these two strategies online

It is a peculiarity of our strategies that

- s_t¹ and s_t² can not be represented as values of a specific gain or loss function. Only general gains can be used.
- the absolute value of one-step gains or losses s_t^i , i = 1, 2, can be unrestrictedly large.

We use the Hannan's Follow the Perturbed Leader FPL algorithm (1957); rediscovered by Kalai and Vempala (2005).

A D > A P > A B > A B >

Prediction with expert advice: general losses

Prediction with expert advice

- s_t^i gain of expert i = 1, ..., N at step $t, s_t^i \in (-\infty, +\infty)$
- s_t Learner gain at step $t, \, s_t \in (-\infty, +\infty)$

$$s_{1:T}^i = \sum_{t=1}^T s_t^i$$
 - cumulative gain of expert *i* on steps $\leq T$

$$m{s}_{1:T} = \sum\limits_{t=1}^{l} m{s}_{t}^{i}$$
 - cumulative gain of Learner on steps $\leq T$

Learner's goal

$$s_{1:T} \ge \max_i s_{1:T}^i - \operatorname{regret}_i$$

・ロ ・ ・ 戸 ・ ・ ヨ ・ ・ ヨ ・

FPL algorithm

Following the Perturbed Leader FPL algorithm:

Output prediction of an expert i which maximizes

$$\mathbf{s}_{1:t-1}^{i}+\frac{1}{\varepsilon}\boldsymbol{\xi}^{i},$$

where ξ^i , i = 1, ..., N, t = 1, 2, ..., is a sequence of i.i.d random variables distributed according to the exponential distribution with the density $p(x) = \exp\{-x\}$, and ε is a learning rate.

Follow Leader algorithm: example

Why randomization? An example.

In the deterministic framework, Learner can perform much worse than each expert:

let the current losses of two experts on steps t = 0, 1, ... be $s_{0,1,2,3,4,5,6,...}^1 = \frac{1}{2}, -1, 1, -1, 1, -1, 1, ...$ and $s_{0,1,2,3,4,5,6,...}^2 = 0, 1, -1, 1, -1, 1, -1, ...$

"Follow Leader" algorithm always chooses the wrong prediction.

The performance of the FPL algorithm

Case: one-step gains of experts are bounded $0 \le s_t^i \le 1$.

In case, where algorithms suffer gains, Kalai and Vempala, Hutter and Poland results can be reformulated such that the expected cumulative gains of the FPL algorithm with variable learning rate $\varepsilon_t = O(1/\sqrt{t})$ has the lower bound

$$E(s_{1:T}) \geq \max_{i=1,\dots,N} s_{1:T}^i - O(\sqrt{T \log N}),$$

ヘロト ヘアト ヘビト ヘビン

where *N* is the number of experts.

Unbounded one-step gains

In that follows we allow gains at any step to be unbounded

 $s_t^i \in (-\infty, +\infty).$

We use only general gains - the notion of a gain function is not used

Loss is a negative gain.



Non-standard scaling

Volume of a game at step t

$$v_t = \sum_{j=1}^t \max_i |s_j^i|.$$

Scaled fluctuation of a game at step t.

$$\operatorname{fluc}(t) = \frac{\Delta v_t}{v_t} = \frac{\max_i |s_t^i|}{v_t},$$

where $\Delta v_t = v_t - v_{t-1}$.

By definition $v_{t-1} \le v_t$ for all t and $0 \le \operatorname{fluc}(t) \le 1$ for all t.

ヘロア 人間 アメヨアメヨア

Volume of the game



Evolution of the volume of the zero-sum game



< ロ > < 回 > < 回 > < 回 > < 回 >

Fluctuation of the game



Fluctuations of the zero-sum game

(日) (四) (三) (三)

Asymptotic consistency

Asymptotic performance of probabilistic algorithms

A probabilistic algorithm is called asymptotically consistent in the mean if

$$\liminf_{T\to\infty}\frac{1}{T}E(s_{1:T}-\max_{i=1,\dots,N}s_{1:T}^{i})\geq 0.$$

A probabilistic algorithm is called asymptotically consistent in the mean in the modified sense if

$$\liminf_{\mathcal{T}\to\infty}\frac{1}{V_{\mathcal{T}}}E(s_{1:\mathcal{T}}-\max_{i=1,\dots,N}s_{1:\mathcal{T}}^{i})\geq 0.$$

A D > A P > A D > A D >

Limits on performance of probabilistic algorithms

Theorem

For any probabilistic algorithm of choosing an expert and for any ε such that $0 < \varepsilon < 1$, two experts exist such that

$$egin{aligned} & v_t o \infty \ as \ t o \infty, \ & ext{fluc}(t) \geq 1 - arepsilon, \ & ext{fluc}(t) \geq 1 - arepsilon, \ & ext{fluc}(t) \geq 1 - arepsilon, \ & ext{fluc}(t) \geq 1 - arepsilon) \end{aligned}$$

for all t, where $s_{1:t}$ is the cumulative gain of this algorithm.

ヘロト ヘアト ヘビト ヘビン

Improving performance

A probabilistic algorithm can be asymptotically consistent only in games where

$$\operatorname{fluc}(t) = \frac{\Delta v_t}{v_t} \to 0 \text{ as } t \to \infty.$$

What is a suitable sufficient condition for a game?

We consider games, such that $\operatorname{fluc}(t) \leq \gamma(t)$ for all t, where $\gamma(t)$ is a computable non-increasing real function such that $0 \leq \gamma(t) \leq 1$ for all t and $\gamma(t) \to 0$ as $t \to \infty$. We consider non-degenerate games, i.e., such that $v_t \to \infty$ as $t \to \infty$.



() > () > () > () > () > ()

FPL algorithm: learning rate

Learning rate of the FPL algorithm

We define a learning rate

$$arepsilon_t = rac{1}{\mu_t v_{t-1}}, ext{ where }$$
 $\mu_t = \sqrt{rac{6}{1+\ln N}} (\gamma(t))^{1/2}.$

<ロ> < □ > < □ > < □ > < □ > < □ >

It holds $\mu_t \leq \mu_{t-1}$ and $v_t \geq v_{t-1}$ for all t.

FPL algorithm

FPL algorithm. FOR t = 1, ..., TDefine

$$I_t = \operatorname{argmax}_{i=1,2,\dots,N} \left\{ s_{1:t-1}^i + \frac{1}{\varepsilon_t} \xi^i \right\},\,$$

where $\varepsilon_t = 1/(\mu_t v_{t-1})$ is the learning rate.

Receive one-step gains s_t^i for experts i = 1, ..., N, and receive one-step gain $s_t = s_t^{l_t}$ of the master algorithm. ENDFOR Here, $\xi^1, ..., \xi^N$ are i.i.d random variables, distributed according to the density $p(x) = \exp\{-x\}$.



FPL algorithm: regret

Theorem

Let a game satisfies

fluc(t) $\leq \gamma(t)$ for all t,

where $\gamma(t)$ is a computable non-increasing real function such that $0 \le \gamma(t) \le 1$ for all *t*.

Then the expected cumulated gain of the FPL algorithm with the variable learning rate is bounded by

$$E(s_{1:T}) \ge \max_{i} s_{1:T}^{i} - 2\sqrt{6(1+\ln N)} \sum_{t=1}^{T} (\gamma(t))^{1/2} \Delta v_{t}.$$



ヘロト ヘロト ヘヨト ヘヨト

FPL algorithm: asymptotic consistency in the mean

Theorem

Let a game satisfies

fluc(t) $\leq \gamma(t)$ for all t,

where $\gamma(t)$ is a computable non-increasing real function such that

- $0 \le \gamma(t) \le 1$ for all t,
- $\gamma(t) \rightarrow 0$, and also,
- $v_t \to \infty$ as $t \to \infty$.

Then the algorithm is asymptotically consistent in the mean

$$\liminf_{T\to\infty}\frac{1}{v_T}E(s_{1:T}-\max_{i=1,\dots,N}s_{1:T}^i)\geq 0.$$

・ロット (雪) ・ (ヨ) ・ (ヨ)

Details of the proof

Details of the proof

We use an auxiliary IFPL algorithm. General scheme:

$$Loss_{1:T}(FPL) \ge Loss_{1:T}(IFPL) - regret_1,$$

$$Loss_{1:T}(IFPL) \ge \max_{i} s_{1:T}^{i} - regret_2,$$

$$Loss_{1:T}(FPL) \ge \max_{i} s_{1:T}^{i} - regret_1 - regret_2$$



IFPL algorithm

IFPL algorithm. FOR t = 1, ..., TDefine the learning rate

$$arepsilon_t' = rac{1}{\mu_t v_t}, ext{ where }$$
 $\mu_t = \sqrt{rac{6}{1+\ln N}} (\gamma(t))^{1/2}$

Choose an expert with the minimal perturbed cumulated loss on steps $\leq t$

$$J_t = \operatorname{argmax}_{i=1,2,\ldots,N} \left\{ \boldsymbol{s}_{1:t}^i + \frac{1}{\varepsilon_t^i} \boldsymbol{\xi}^i \right\}.$$

Receive the one step loss $\tilde{s}_t = s_t^{J_t}$ of the IFPL algorithm. ENDFOR



Lemma 1

Lemma

The cumulated expected losses of the FPL and IFPL algorithms satisfy the inequality

$$E(s_{1:T}) \ge E(\widetilde{s}_{1:T}) - 6\sum_{t=1}^{T} (\gamma(t))^{1-lpha_t} \Delta v_t$$

for all *T*, where $\alpha_t < 1$.

The optimal choice

$$\alpha_t = \frac{1}{2} \left(1 - \frac{\ln \frac{1 + \ln N}{6}}{\ln \gamma(t)} \right)$$



(ロ) (四) (日) (日) (日)

Lemma 2

Lemma

The expected cumulative loss of the IFPL algorithm is bounded

$$E(\tilde{s}_{1:T}) \geq \max_{i} s_{1:T}^{i} - (1 + \ln N) \sum_{t=1}^{T} (\gamma(t))^{\alpha_{t}} \Delta v_{t}$$

for all T.

By definition $\mu_t = (\gamma(t))^{\alpha_t}$.



Polynomially bounded experts

Corollary

Assume that $|s_t^i| \le t^a$ for all t and i = 1, ..., N, and

$$\liminf_{t\to\infty}\frac{v_t}{t^{a+\delta}}>0,$$

where a and δ are positive real numbers. Let also, $\gamma(t) = t^{-\delta}$ and μ_t is defined above. Then

- our algorithm is asymptotically consistent in the mean for any a > 0 and δ > 0;
- the expected cumulated gain of this algorithm is bounded

$$E(s_{1:T}) \geq \max_{i} s_{1:T}^{i} - O\left(\sqrt{\ln N}T^{1-\frac{1}{2}\delta+a}\right)$$

as $T \rightarrow \infty$.

Derandomization

Derandomized version of the algorithm

Learner's gain at step t

$$G_t = P\{I_t = 1\}s_t^1 + P\{I_t = 2\}s_t^2$$

Learner's cumulative gain at steps $\leq T$

$$G_{1:T} = \sum_{t=1}^{T} G_t = E(s_{1:T})$$

Lower bound

$$G_{1:T} \ge |s_{1:T}^1| - 8\sum_{t=1}^T (\gamma(t))^{1/2} |s_t^1|$$

・ロト ・四ト ・ヨト ・ヨト

Derandomization

Corollary

The FPL algorithm is asymptotically riskless

$$\liminf_{T\to\infty}\frac{1}{v_T}G_{1:T}\geq 0.$$



Merging strategies



Gain of the FL (Follow Leader) algorithm (without randomization). Green and Blue lines - gains of experts; Red line - Learner gain

Merging strategies



The FPL algorithm's expected gain.

Green and Blue lines - gains of experts; Red line - Learner gain

Merging strategies



Some runs of FPL are excellent

・ロト ・ 日 ・ ・ 回 ・ ・ 回

Conclusions

Conclusions

We have studied two different problems:

- How to use the fractional Brownian motion of prices to suffer gain with "minimal" risk on financial market;
- How to extend methods of the theory of prediction with expert advice for the case when experts one-step gains are unbounded.

We use a solution of the second problem to solve the first one.



Stock price



Evolution of the stock price (s=1000,1000)

・ロ・・ 日本・ ・ 日本・ ・ 日本

Gain from high volatility. Constant leader



Gain of the FL algorithm (s=1000,1000)



・ロ・・ 日・ ・ 回・ ・

Gain from high volatility. Constant leader



Expected gain of the FPL algorithm (s=1000,1000)



• • • • • • • • • • • •

Stock price



Evolution of the stock price (s1=4100,1000)



・ロ・・ 日・ ・ 回・ ・

Changing the leader



Gain of the FL algorithm (s1=4100,1000)



ъ

・ロ・・ 日本・ ・ 日本・

Changing the leader



Expected gain of the FPL algorithm (s1=4100,1000)

